

Vector Reachability Problem in $SL(2, \mathbb{Z})$ [☆]

Igor Potapov^a, Pavel Semukhin^{a,*}

^a*Department of Computer Science, University of Liverpool, United Kingdom*

Abstract

This paper solves two open problems about the decidability of the vector reachability problem over a finitely generated semigroup of matrices from $SL(2, \mathbb{Z})$ and the point to point reachability (over rational numbers) for fractional linear transformations, where associated matrices are from $SL(2, \mathbb{Z})$. The approach to solving reachability problems is based on the characterization of reachability paths between points which is followed by the translation of numerical problems on matrices into computational and combinatorial problems on words and formal languages. We will also use this technique to prove that a special case of the scalar reachability problem is decidable.

Keywords: Matrix Semigroup, Vector Reachability Problem, Special Linear Group, Linear Fractional Transformation, Automata and Formal Languages
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1. Introduction

Decision problems on matrices were intensively studied from 1947 when A. Markov showed the connection between classical computations and problems for matrix semigroups [1]. Moreover matrix products play an essential role
5 in the representation of various computational processes, i.e., linear recurrent sequences [2, 3, 4], arithmetic circuits [5], hybrid and dynamical systems [6, 7], probabilistic and quantum automata [8], stochastic games, broadcast protocols [9], optical systems, etc. New algorithms for solving reachability problems in
10 matrix semigroups can be incorporated into software verification tools and used for analysis of mathematical models in physics, chemistry, biology, ecology, and economics.

However, many computational problems for matrix semigroups are inherently difficult to solve even when the problems are considered in dimension two, and most of these problems become undecidable in general starting from

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*Corresponding author

Email addresses: potapov@liverpool.ac.uk (Igor Potapov), semukhin@liverpool.ac.uk (Pavel Semukhin)

15 dimension three or four. Examples of such problems are the Membership problem (including the special cases of the Mortality and Identity problems), vector reachability, scalar reachability, freeness problem and the emptiness problem of matrix semigroups intersection [10]. All above problems are tightly connected, including three central problems:

- 20 • **The membership problem:** Let $S = \langle G \rangle$ be a semigroup generated by a finite set G of $n \times n$ matrices. Determine whether a given matrix M belongs to S , that is, determine whether there exists a sequence of matrices M_1, M_2, \dots, M_k in G such that $M = M_1 \cdot M_2 \cdot \dots \cdot M_k$
- 25 • **The vector reachability problem:** Let \mathbf{x} and \mathbf{y} be two vectors and S be a given finitely generated semigroup of $n \times n$ matrices. Determine whether there is a matrix $M \in S$ such that $M\mathbf{x} = \mathbf{y}$.
- **The scalar reachability problem:** Let \mathbf{x} and \mathbf{y} be two vectors, λ be a scalar, and S be a given finitely generated semigroup of $n \times n$ matrices. Determine whether there is a matrix $M \in S$ such that $\mathbf{x}^\top M \mathbf{y} = \lambda$.

30 The vector reachability problem can be seen as a parameterized version of the membership problem, where some elements of a matrix M are either independent variables or variables linked by some equations. In contrast to the original membership problem, where all values of M are constants, in vector reachability we may have an infinite set of matrices that can transform a vector \mathbf{x} to \mathbf{y} . Thus the decidability results for the membership cannot be directly applied to the vector reachability problem.

The scalar reachability can be viewed as a vector to hyperplane reachability problem. Indeed, we can rewrite the equation $\mathbf{x}^\top M \mathbf{y} = \lambda$ as a system of two equations: $M\mathbf{y} = \mathbf{z}$ and $\mathbf{x}^\top \mathbf{z} = \lambda$. So, the question becomes if there is a matrix $M \in S$ that maps a given vector \mathbf{y} to a vector \mathbf{z} that lies on a hyperplane $\mathbf{x}^\top \mathbf{z} = \lambda$. Because there are infinitely many vectors on a hyperplane, decidability of the scalar reachability problem does not follow directly from the decidability of the vector reachability problem.

Most of the problems such as membership, vector reachability and freeness are undecidable for 3×3 integer matrices. The undecidability proofs in matrix semigroups are mainly based on various techniques and methods of embedding universal computations into three and four dimensional matrices and their products. The case of dimension two is the most intriguing one since there is some evidence that if these problems are undecidable, then this cannot be proved using a construction similar to the one used for dimensions 3 and 4. In particular, there is no injective semigroup morphism from pairs of words over any finite alphabet (with at least two elements) into complex 2×2 matrices [11], which means that the coding of independent pairs of words in 2×2 complex matrices is impossible and the exact encoding of the Post Correspondence Problem or a computation of a Turing Machine cannot be used directly for proving undecidability in 2×2 matrix semigroups over \mathbb{Z} , \mathbb{Q} or \mathbb{C} . The only undecidability result

in dimension two for the vector reachability and the membership problems has been shown in the case of 2×2 matrices over quaternions [12].

The main hypothesis is that problems for 2×2 matrix semigroups over integers, rationals or complex numbers could be decidable, but not much is known about the status of these problems. There was some progress on the *Membership problem*, which was shown to be decidable in $\text{SL}(2, \mathbb{Z})$, and the *Identity problem*, which was shown to be decidable in $\mathbb{Z}^{2 \times 2}$ [13]. Later the decidability of the *Freeness problem* (that is, to decide whether each element can be expressed uniquely as a product of generating matrices) was shown for $\text{SL}(2, \mathbb{Z})$ [14] and for upper-triangular 2×2 matrices with rational entries when the products are restricted to certain bounded languages [15]. The Mortality, Identity and Vector Reachability problems were shown to be at least NP-hard for $\text{SL}(2, \mathbb{Z})$ [16, 10], but for the modular group the Membership was shown to be decidable in polynomial time by Gurevich and Schupp [17].

Recently, the Membership problem for nonsingular matrices from $\mathbb{Z}^{2 \times 2}$ was proven to be decidable [18]. Furthermore, it was shown that the following problems in $\text{SL}(2, \mathbb{Z})$ are NP-complete: Identity, Membership [19] and Non-Freeness [19, 20].

The algorithmic properties of $\text{SL}(2, \mathbb{Z})$ are important in the context of many fundamental problems in hyperbolic geometry [21, 22, 23], dynamical systems [24], Lorenz/modular knots [25], braid groups [26], particle physics, high energy physics [27], M/string theories [28], ray tracing analysis, music theory [29] and can lead to further decidability results in $\mathbb{Z}^{2 \times 2}$ using matrix presentation in the Smith normal form.

This paper solves two open problems about the decidability of the vector reachability problem for finitely generated semigroups of matrices from $\text{SL}(2, \mathbb{Z})$ and the point to point reachability (over rational numbers) for fractional linear transformations $f_M(x) = \frac{ax+b}{cx+d}$, where the associated matrix $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$ belongs to $\text{SL}(2, \mathbb{Z})$. The approach to solving these reachability problems for 2×2 matrix semigroups is based on the analysis of reachability paths between vectors or points. This analysis is then used to translate the numerical reachability problems into computational problems on words and regular languages. We also present several extensions of our main results, give a geometric interpretation of reachability paths, and use this technique to solve a special case of the scalar reachability problem.

The decidability proof of the vector reachability problem in dimension two presented in this paper is the first nontrivial new result for solving vector reachability problems since 1996 when it was shown that the problem is decidable for any commutative matrix semigroup in any dimension [30] and for a special case of non-commuting matrices [31]. On the other hand, in the general case of non-commuting matrices the problem is known to be undecidable already for integer matrices in dimension three [32].

2. Preliminaries

100 The integers and rationals are denoted by \mathbb{Z} and \mathbb{Q} , respectively, and $\text{SL}(2, \mathbb{Z})$ is a group of 2×2 integer matrices with determinant 1. The notation $a \mid b$ means that a divides b , and $a \nmid b$ means that a does not divide b , when a and b are integer numbers.

Definition 1. With each matrix $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix} \in \text{SL}(2, \mathbb{Z})$ we associate a fractional linear map (also called Möbius transformation) $f_M : \mathbb{Q} \rightarrow \mathbb{Q}$ defined as

$$f_M(x) = \frac{ax + b}{cx + d}.$$

This definition can be extended to $f : \mathbb{Q} \cup \{\infty\} \rightarrow \mathbb{Q} \cup \{\infty\}$ in a natural way by setting $f_M(\infty) = \frac{a}{c}$ if $c \neq 0$, $f_M(\infty) = \infty$ if $c = 0$, and $f_M(x) = \infty$ if $cx + d = 0$.

Note that we have $f_{M_1} \circ f_{M_2} = f_{M_1 M_2}$ for any matrices M_1 and M_2 .

110 Let M_1, \dots, M_n be a finite collection of matrices. Then $\langle M_1, \dots, M_n \rangle$ denotes the multiplicative semigroup (including the identity matrix) generated by M_1, \dots, M_n .

Definition 2. The *vector reachability problem (VRP)* in $\text{SL}(2, \mathbb{Z})$ is defined as follows: Given two vectors \mathbf{x} and \mathbf{y} with integer coefficients and a finite collection of matrices M_1, \dots, M_n from $\text{SL}(2, \mathbb{Z})$, decide whether there exists a matrix $M \in \langle M_1, \dots, M_n \rangle$ such that $M\mathbf{x} = \mathbf{y}$.

Definition 3. The *reachability problem by fractional linear transformations (FLT)* in $\text{SL}(2, \mathbb{Z})$ is defined as follows: Given two rational numbers x and y and a finite collection of matrices M_1, \dots, M_n from $\text{SL}(2, \mathbb{Z})$, decide whether there exists a matrix $M \in \langle M_1, \dots, M_n \rangle$ such that $f_M(x) = y$.

3. Overview of the main results

The main result of our paper is that the vector reachability problem and the reachability problem by fractional linear transformations for $\text{SL}(2, \mathbb{Z})$ are decidable (Theorem 16). Both proofs follow the same pattern. We will use the fact that any matrix M from $\text{SL}(2, \mathbb{Z})$ can be expressed as product of matrices $S = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}$ and $R = \begin{bmatrix} 0 & -1 \\ 1 & 1 \end{bmatrix}$. So we can represent any $M \in \text{SL}(2, \mathbb{Z})$ by a word w in the alphabet $\Sigma = \{S, R\}$.

The main idea of the proof is to show that the solution set of the equation $M\mathbf{x} = \mathbf{y}$ has the form

$$\left\{ B \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}^t C : t \in \mathbb{Z} \right\},$$

130 where B and C are some matrices from $\text{SL}(2, \mathbb{Z})$ that can be computed in PTIME from \mathbf{x}, \mathbf{y} (Theorem 10). Similarly, the solution set of the equation $f_M(x) = y$ can be presented as a union of two sets of such form (Theorem 12).

After translating matrices into words, these sets become regular languages. On the other hand, the language that corresponds to the semigroup $\langle M_1, \dots, M_n \rangle$ is also regular. Indeed, if M_i corresponds to the word w_i , then the semigroup $\langle M_1, \dots, M_n \rangle$ translates into the language $(w_1 + \dots + w_n)^*$. The last step of the proof is to show that the emptiness problem of the intersection of two such languages is decidable (Proposition 15).

Here is a more detailed description of our proofs. Let $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$, $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$ and $\mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$. To show that the equation $M\mathbf{x} = \mathbf{y}$ defines a regular language we must solve the following system of three equations in four unknown variables:

$$x_1a + x_2b = y_1 \quad x_1c + x_2d = y_2 \quad ad - bc = 1$$

Choosing b as a free parameter, we can reduce it to the following system of linear congruence equations:

$$\begin{aligned} x_2b &\equiv y_1 \pmod{x_1} \\ y_2b &\equiv -x_1 \pmod{y_1} \\ x_2y_2b &\equiv y_1y_2 - x_1x_2 \pmod{x_1y_1} \end{aligned}$$

By Lemma 6 from Section 4, the above system either has no solutions or it has a solution of the form $b = b_1t + b_2$, where $t \in \mathbb{Z}$, and hence all coefficients of the matrix M are linear functions of t . In Proposition 9 we will show that such matrices can be written in the form

$$M = B \begin{bmatrix} 1 & k \\ 0 & 1 \end{bmatrix}^t C,$$

where B, C are some matrices from $\text{SL}(2, \mathbb{Z})$, k is a fixed integer number and $t \in \mathbb{Z}$ is a free parameter. After that it is not hard to see that such solution translates into a regular language.

We will use a similar approach to prove that the equation $f_M(x) = y$ also defines a regular language. In fact, we will do it by showing that the solution set of $f_M(x) = y$ is equal to the union of the solution sets of the equations $M\mathbf{x} = \mathbf{y}$ and $M\mathbf{x} = -\mathbf{y}$ for suitable vectors \mathbf{x} and \mathbf{y} .

The final step is to show that there is an algorithm that decides whether the intersection of two regular subsets of $\text{SL}(2, \mathbb{Z})$ is empty or not. Our idea relies on the fact that the intersection of two regular languages is regular, and that the emptiness problem for regular languages is decidable. The problem here is that we cannot apply these facts directly because for each matrix $M \in \text{SL}(2, \mathbb{Z})$ there are infinitely many words $w \in \{S, R\}^*$ that correspond to M , and only some of them may appear in the given language. However there is only one *reduced* word that corresponds to M , that is, the word that does not have a substring of the form SS or RRR . So, our solution is to take any automaton \mathcal{A} and turn it into a new automaton $\tilde{\mathcal{A}}$ that accepts the same language as \mathcal{A} plus all reduced words w that correspond to non-reduced words w' accepted by \mathcal{A} .

The construction of the automaton $\tilde{\mathcal{A}}$ was inspired by a similar construction from [13]. Note that in $\text{SL}(2, \mathbb{Z})$ we have an equality $S^2 = R^3 = -I$. Thus to construct $\tilde{\mathcal{A}}$ we add to \mathcal{A} a new ε -transition from a state q_1 to a state q_2 if there is a run of \mathcal{A} from q_1 to q_2 labelled by SS or RRR . We will apply this procedure iteratively until no new ε -transitions can be added. However we need to keep track of sign changes when we add new ε -transitions. To achieve this we will use *signed automata*, which are slight modifications of the usual finite automata but they take into account such sign changes.

Now to solve the emptiness problem for the intersection of two regular subsets of $\text{SL}(2, \mathbb{Z})$, which are defined by regular languages L_1 and L_2 , we take the signed automata \mathcal{A}_1 and \mathcal{A}_2 that accept L_1 and L_2 , respectively, and construct new automata $\tilde{\mathcal{A}}_1$ and $\tilde{\mathcal{A}}_2$ as described above. After that we can check whether $L(\tilde{\mathcal{A}}_1) \cap L(\tilde{\mathcal{A}}_2) \neq \emptyset$.

In the end of Section 5 we will show how to extend these decidability results to arbitrary regular subsets of $\text{SL}(2, \mathbb{Z})$, i.e., subsets that are defined by finite automata. Using this technique we will show how to algorithmically solve the equation

$$M_1^{x_1} \cdots M_k^{x_k} \mathbf{x} = N_1^{y_1} \cdots N_l^{y_l} \mathbf{y},$$

where \mathbf{x}, \mathbf{y} are given vectors from $\mathbb{Z} \times \mathbb{Z}$, the matrices M_1, \dots, M_k and N_1, \dots, N_l are from $\text{SL}(2, \mathbb{Z})$, and x_1, \dots, x_k and y_1, \dots, y_l are unknown non-negative integers.

Furthermore, in Section 6 we will consider geometric properties of the solutions of the equations $M\mathbf{x} = \mathbf{y}$ and $f_M(x) = y$ and show how to apply them to prove that a special case of the scalar reachability problem is decidable.

4. Characterization of the matrices that satisfy the equations $M\mathbf{x} = \mathbf{y}$ and $f_M(x) = y$

The characterization of the solution set of the equation $M\mathbf{x} = \mathbf{y}$ given in Theorem 10 will follow from Propositions 7 and 9. First, we prove one simple lemma which we will use several times in our arguments.

Lemma 4. Let $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$ and $\mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$ be vectors from $\mathbb{Z} \times \mathbb{Z}$ and M be a matrix from $\text{SL}(2, \mathbb{Z})$ such that $M\mathbf{x} = \mathbf{y}$. Then $\gcd(x_1, x_2) = \gcd(y_1, y_2)$.

Proof. Take any $k \in \mathbb{Z}$ such that $k \mid x_1, x_2$ and let $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$. Then from $M\mathbf{x} = \mathbf{y}$ we have $y_1 = ax_1 + bx_2$ and $y_2 = cx_1 + dx_2$. Thus $k \mid y_1, y_2$. Now since $M \in \text{SL}(2, \mathbb{Z})$, M^{-1} is also in $\text{SL}(2, \mathbb{Z})$, and $M\mathbf{x} = \mathbf{y}$ is equivalent to $M^{-1}\mathbf{y} = \mathbf{x}$. So, if $k \in \mathbb{Z}$ is any number such that $k \mid y_1, y_2$, then $k \mid x_1, x_2$. Therefore, $\gcd(x_1, x_2) = \gcd(y_1, y_2)$. \square

For the proof of Proposition 7 we will need the following two lemmas.

Lemma 5. *Consider a linear congruence equation*

$$ax \equiv b \pmod{n}.$$

If $\gcd(a, n) \nmid b$, then the equation has no solution. If $\gcd(a, n) \mid b$, then all solutions of the equation can be written in the form $x \equiv c \pmod{\frac{n}{\gcd(a, n)}}$ for some c . Moreover, there is a polynomial time algorithm that determines whether such an equation has a solution and if so, finds it.

Proof. Given a and n , using Euclidean algorithm we can find in polynomial time $d = \gcd(a, n)$ and integer numbers u and v such that $d = ua + vn$. Equation $ax \equiv b \pmod{n}$ can be written as $ax = b + kn$, where $k \in \mathbb{Z}$. It is clear that if $d \nmid b$, then there is no solution. Otherwise, let $b = b'd$, $a = a'd$, and $n = n'd$. Then our equation is equivalent to $a'x \equiv b' \pmod{n'}$. Furthermore, we have $ua' + vn' = 1$ and hence $ua' \equiv 1 \pmod{n'}$. Thus

$$x \equiv (ua')x \equiv u(a'x) \equiv ub' \pmod{n'}.$$

Note that all these computations can be done in polynomial time. \square

Lemma 6. *Consider a system of two linear congruence equations*

$$\begin{aligned} a_1x &\equiv b_1 \pmod{n_1} \\ a_2x &\equiv b_2 \pmod{n_2} \end{aligned} \tag{1}$$

Then either such system has no solution, or all its solutions are of the form $x \equiv c \pmod{n}$ for some c and $n \mid n_1n_2$. Moreover, there is a polynomial time algorithm that determines whether (1) has a solution and if so, finds it.

Proof. Using the algorithm of Lemma 5, we can solve each equation separately. If one of them does not have a solution, then the system (1) also has no solution. Suppose the first and second equation have the solutions

$$x \equiv c_1 \pmod{n'_1} \quad \text{and} \quad x \equiv c_2 \pmod{n'_2},$$

respectively, which can be found in PTIME. Note that $n'_i \mid n_i$ for $i = 1, 2$.

Let $n = \text{lcm}(n'_1, n'_2)$. We can rewrite the solutions as

$$\begin{aligned} x &\equiv c_1, c_1 + n'_1, c_1 + 2n'_1, \dots, c_1 + (n''_1 - 1)n'_1 \pmod{n}, \\ x &\equiv c_2, c_2 + n'_2, c_2 + 2n'_2, \dots, c_2 + (n''_2 - 1)n'_2 \pmod{n}, \end{aligned}$$

where $n''_1 = n/n'_1$ and $n''_2 = n/n'_2$. Let

$$\begin{aligned} A_1 &= \{c_1, c_1 + n'_1, c_1 + 2n'_1, \dots, c_1 + (n''_1 - 1)n'_1\}, \\ A_2 &= \{c_2, c_2 + n'_2, c_2 + 2n'_2, \dots, c_2 + (n''_2 - 1)n'_2\}. \end{aligned}$$

Note that $A_1 \cap A_2$ contains at most one element. Indeed, if $c, c' \in A_1 \cap A_2$, then $n'_1 \mid c - c'$ and $n'_2 \mid c - c'$. Hence $n = \text{lcm}(n'_1, n'_2) \mid c - c'$. Since $|c - c'| < n$, we have $c = c'$.

Now if $A_1 \cap A_2$ is empty, then (1) has no solution. If $A_1 \cap A_2 = \{c\}$, then the solution of (1) is $x \equiv c \pmod{n}$. To find this solution in PTIME, observe the following. The equations $x \equiv c_1 \pmod{n'_1}$ and $x \equiv c_2 \pmod{n'_2}$ are equivalent to $x = c_1 + kn'_1$ and $x = c_2 + ln'_2$, respectively, where $k, l \in \mathbb{Z}$. To find the intersection of these solutions we set $c_1 + kn'_1 = c_2 + ln'_2$, which is equivalent to $c_1 - c_2 = ln'_2 - kn'_1$. Using Euclidean algorithm, we can find in PTIME $d = \gcd(n'_1, n'_2)$ and integer numbers u, v such that

$$d = un'_1 + vn'_2. \quad (2)$$

Obviously, if $d \nmid c_1 - c_2$, then there is no solution. So suppose $c_1 - c_2 = hd$, for some $h \in \mathbb{Z}$. Multiplying (2) by h we obtain

$$c_1 - c_2 = hd = (hu)n'_1 + (hv)n'_2 \quad \text{or} \quad c_1 - (hu)n'_1 = c_2 + (hv)n'_2.$$

Let c be the number in the set $\{0, \dots, n-1\}$ such that

$$c \equiv c_1 - (hu)n'_1 = c_2 + (hv)n'_2 \pmod{n}.$$

Then $x \equiv c \pmod{n}$ is the desired solution. It is not hard to see that the above algorithm runs in polynomial time. \square

Proposition 7. Let $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$ and $\mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$ be two vectors from $\mathbb{Z} \times \mathbb{Z}$, such that \mathbf{x} is not equal to the zero vector $\mathbf{0}$, and consider the matrix equation $M\mathbf{x} = \mathbf{y}$, where M is an unknown matrix from $\text{SL}(2, \mathbb{Z})$. Then either this equation does not have a solution or all its solutions are given by

$$M = tA_1 + A_2,$$

where t is any integer number, A_1, A_2 are some matrices from $\mathbb{Z}^{2 \times 2}$ such that A_1 is a nonzero matrix. Moreover, there is a polynomial time algorithm that determines whether such an equation has a solution and if so, finds it.

Proof. Let $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$ be a matrix that satisfies the equations $M\mathbf{x} = \mathbf{y}$ and $\det(M) = 1$. So we have the following system of equations:

$$\begin{aligned} x_1a + x_2b &= y_1 \\ x_1c + x_2d &= y_2 \\ ad - bc &= 1 \end{aligned} \quad (3)$$

Recall that by assumption $\mathbf{x} \neq \mathbf{0}$. Without loss of generality, suppose that $x_1 \neq 0$. In this case we have

$$a = \frac{y_1 - x_2b}{x_1} \quad \text{and} \quad c = \frac{y_2 - x_2d}{x_1}.$$

Substituting these values for a and c into the equation $ad - bc = 1$, we obtain

$$(y_1 - x_2b)d - (y_2 - x_2d)b = x_1$$

or, equivalently, $y_1d - y_2b = x_1$. If $y_1 = y_2 = 0$, then there is no solution because by assumption $x_1 \neq 0$. Again, without loss of generality, assume that $y_1 \neq 0$. Hence $d = \frac{x_1 + y_2b}{y_1}$. If we choose b as a free parameter, then the general solution of the system of equations (3) will be:

$$\begin{aligned} a &= \frac{y_1 - x_2b}{x_1}, & d &= \frac{x_1 + y_2b}{y_1}, \\ c &= \frac{y_2 - x_2 \frac{x_1 + y_2b}{y_1}}{x_1} = \frac{y_1y_2 - x_1x_2 - x_2y_2b}{x_1y_1}. \end{aligned}$$

We are interested only in integer solutions, that is, when a , c , and b are in \mathbb{Z} , which means that b must satisfy the following congruences:

$$\begin{aligned} x_2b &\equiv y_1 \pmod{x_1} \\ y_2b &\equiv -x_1 \pmod{y_1} \\ x_2y_2b &\equiv y_1y_2 - x_1x_2 \pmod{x_1y_1} \end{aligned}$$

250 Applying the algorithm from Lemma 6 two times, we can determine in PTIME whether the above system has a solution or not. If the solution exists, the algorithm outputs it in the form $b \equiv b_2 \pmod{b_1}$, where $b_1 \mid x_1y_1$.

So, the coefficient b is of the form $b = b_1t + b_2$, where $t \in \mathbb{Z}$. Substituting this expression for b in the formulas for a , c , and d we obtain:

$$\begin{aligned} a &= \frac{y_1 - x_2b_2 - x_2b_1t}{x_1} = a_1t + a_2, & d &= \frac{x_1 + y_2b_2 + y_2b_1t}{y_1} = d_1t + d_2, \\ c &= \frac{y_1y_2 - x_1x_2 - x_2y_2b_2 - x_2y_2b_1t}{x_1y_1} = c_1t + c_2, \end{aligned}$$

255 where a_i , c_i , and d_i , for $i = 1, 2$, are some constants which are necessarily in \mathbb{Z} because if we let $t = 0$ or $t = 1$ in the above expressions they must evaluate to integer numbers. Therefore, the solution to the system of equations (3) can be written as:

$$M = \begin{bmatrix} a_1t + a_2 & b_1t + b_2 \\ c_1t + c_2 & d_1t + d_2 \end{bmatrix} = t \begin{bmatrix} a_1 & b_1 \\ c_1 & d_1 \end{bmatrix} + \begin{bmatrix} a_2 & b_2 \\ c_2 & d_2 \end{bmatrix},$$

where t is any integer number. To complete the proof we set $A_1 = \begin{bmatrix} a_1 & b_1 \\ c_1 & d_1 \end{bmatrix}$

260 and $A_2 = \begin{bmatrix} a_2 & b_2 \\ c_2 & d_2 \end{bmatrix}$. Note that A_1 is a nonzero matrix since at least one of its coefficients, namely b_1 , is not equal to zero. Furthermore, the above algorithm runs in polynomial time because the only nontrivial step is to solve the system of linear congruence equations, which according to Lemma 6 can be done in PTIME. \square

265 For the next proposition we will need the following theorem about the Smith normal form of a matrix.

Theorem 8 (Smith normal form [33]). *For any nonzero matrix $A \in \mathbb{Z}^{2 \times 2}$, there are matrices B, C from $\text{SL}(2, \mathbb{Z})$ such that*

$$A = B \begin{bmatrix} t_1 & 0 \\ 0 & t_2 \end{bmatrix} C$$

for some $t_1, t_2 \in \mathbb{Z}$ such that $t_1 \neq 0$ and $t_1 \mid t_2$. Moreover, B, C, t_1, t_2 can be computed in polynomial time.

Proposition 9. *Let A_1 and A_2 be matrices from $\mathbb{Z}^{2 \times 2}$ such that A_1 is a nonzero matrix and, for every $t \in \mathbb{Z}$, we have $tA_1 + A_2 \in \text{SL}(2, \mathbb{Z})$. Then there are matrices B and C from $\text{SL}(2, \mathbb{Z})$ and $k \in \mathbb{Z}$ such that*

$$tA_1 + A_2 = BT^{kt}C \text{ for every } t \in \mathbb{Z},$$

where $T = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} \in \text{SL}(2, \mathbb{Z})$. Moreover, B, C , and k can be computed in polynomial time.

Proof. Let $A_1 = \begin{bmatrix} a_1 & b_1 \\ c_1 & d_1 \end{bmatrix}$ and $A_2 = \begin{bmatrix} a_2 & b_2 \\ c_2 & d_2 \end{bmatrix}$. By the assumption, for every $t \in \mathbb{Z}$, we have $\begin{vmatrix} a_1t + a_2 & b_1t + b_2 \\ c_1t + c_2 & d_1t + d_2 \end{vmatrix} = 1$. That is

$$(a_1t + a_2)(d_1t + d_2) - (b_1t + b_2)(c_1t + c_2) = 1 \text{ or} \\ (a_1d_1 - b_1c_1)t^2 + (a_1d_2 + a_2d_1 - b_1c_2 - b_2c_1)t + a_2d_2 - b_2c_2 = 1 \text{ for all } t \in \mathbb{Z}.$$

Therefore, $a_1d_1 - b_1c_1 = 0$, $a_1d_2 + a_2d_1 - b_1c_2 - b_2c_1 = 0$, and $a_2d_2 - b_2c_2 = 1$. In particular, $\det(A_1) = 0$ and $\det(A_2) = 1$.

By Theorem 8, there are matrices F and G from $\text{SL}(2, \mathbb{Z})$ such that

$$A_1 = F \begin{bmatrix} k & 0 \\ 0 & l \end{bmatrix} G$$

for some $k, l \in \mathbb{Z}$ such that $k \mid l$. Since $\det(A_1) = 0$ we have that $kl = 0$. However if $k = 0$ and $l = 0$, then A_1 is equal to the zero matrix, contrary to the assumption. Hence we must have that $k \neq 0$ and $l = 0$.

Now $F^{-1}(tA_1 + A_2)G^{-1} = \begin{bmatrix} kt + a & b \\ c & d \end{bmatrix}$, for some $a, b, c, d \in \mathbb{Z}$. Note that since $\det(F) = \det(G) = \det(tA_1 + A_2) = 1$, we have

$$\begin{vmatrix} kt + a & b \\ c & d \end{vmatrix} = dkt + ad - bc = 1 \text{ for every } t \in \mathbb{Z}.$$

Hence $dk = 0$ and so $d = 0$. Substituting $d = 0$ in the above equation, we obtain $bc = -1$. Since b and c are integers, there are only two possibilities: $b = 1, c = -1$, or $b = -1, c = 1$. So the above matrix actually looks like

$$F^{-1}(tA_1 + A_2)G^{-1} = \begin{bmatrix} kt + a & \mp 1 \\ \pm 1 & 0 \end{bmatrix}.$$

Therefore,

$$T^{-c(kt+a)}F^{-1}(tA_1 + A_2)G^{-1} = D,$$

where $c = \pm 1$ and $D = \begin{bmatrix} 0 & \mp 1 \\ \pm 1 & 0 \end{bmatrix} \in \text{SL}(2, \mathbb{Z})$. Hence

$$tA_1 + A_2 = FT^{(ck)t}T^{ca}DG.$$

Note that F and $T^{ca}DG$ are in $\text{SL}(2, \mathbb{Z})$. This completes the proof. The bound on complexity follows from the fact that F and G can be computed in PTIME by Theorem 8. \square

As a corollary of Propositions 7 and 9 we obtain the following theorem.

Theorem 10. Let $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$ and $\mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$ be vectors from $\mathbb{Z} \times \mathbb{Z}$ such that $\mathbf{x} \neq \mathbf{0}$, and consider the matrix equation $M\mathbf{x} = \mathbf{y}$, where M is an unknown matrix from $\text{SL}(2, \mathbb{Z})$. Then either this equation does not have a solution or all its solutions are given by the following formula

$$M = B \begin{bmatrix} 1 & k \\ 0 & 1 \end{bmatrix}^t C, \text{ where } t \in \mathbb{Z}.$$

In the above expression B and C are some matrices from $\text{SL}(2, \mathbb{Z})$, and k is an integer number. Moreover, there is a polynomial time algorithm that determines whether such an equation has a solution and if so, finds the suitable matrices B , C and the integer k .

In Section 6 we will give a geometric interpretation of reachability paths (Figure 2 and Proposition 18), using which we can prove the following corollary.¹

Corollary 11. The value of the parameter k in Theorem 10 is equal to 1.

Proof. Our proof relies on Proposition 18 from Section 6. We need to show that

$$\{BT^{kt}C : t \in \mathbb{Z}\} = \{BT^tC : t \in \mathbb{Z}\}.$$

Note that the inclusion $\{BT^{kt}C : t \in \mathbb{Z}\} \subseteq \{BT^tC : t \in \mathbb{Z}\}$ is obvious. On the other hand, by Proposition 18, the vectors $\mathbf{u} = C\mathbf{x}$ and $\mathbf{v} = B^{-1}\mathbf{y}$ are equal to each other and have the form $\begin{bmatrix} d \\ 0 \end{bmatrix}$. Since $T = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}$, it is easy to check that $T\mathbf{u} = \mathbf{u} = \mathbf{v}$. From this we can conclude that $T^t\mathbf{u} = \mathbf{v}$ and hence

$$BT^tC\mathbf{x} = \mathbf{y} \text{ for any } t \in \mathbb{Z}.$$

¹Even though we use Corollary 11 in the proofs of Theorem 12 and Proposition 20, it is not essential there for proving decidability. Namely, all references to Corollary 11 in these proofs can be replaced by references to Theorem 10, at the same time replacing T with T^k where appropriate.

In other words, matrices of the form BT^tC transform \mathbf{x} to \mathbf{y} . However, in
 310 Theorem 10 we proved that all such matrices belong to the set $\{BT^{kt}C : t \in \mathbb{Z}\}$.
 Therefore, we obtain the inclusion $\{BT^tC : t \in \mathbb{Z}\} \subseteq \{BT^{kt}C : t \in \mathbb{Z}\}$, and this
 proves the corollary. \square

Theorem 10 provides us with a characterization of the matrices $M \in \text{SL}(2, \mathbb{Z})$
 that map vector \mathbf{x} to vector \mathbf{y} . This characterization will be used later to
 315 prove the decidability of the vector reachability problem. We now give a similar
 characterization of the matrices $M \in \text{SL}(2, \mathbb{Z})$ for which the fractional linear
 transformation f_M maps a number x to number y . In fact, we will do this by
 reducing the problem to finding the solutions of the equation $M\mathbf{x} = \mathbf{y}$ which
 we discussed above.

320 **Theorem 12.** *Let x and y be rational numbers and let $\mathcal{F}(x, y)$ be the following
 set of matrices from $\text{SL}(2, \mathbb{Z})$:*

$$\mathcal{F}(x, y) = \{M \in \text{SL}(2, \mathbb{Z}) : f_M(x) = y\}.$$

Then $\mathcal{F}(x, y) = \mathcal{F}_1(x, y) \cup \mathcal{F}_2(x, y)$, where each $\mathcal{F}_i(x, y)$ is either empty or has
 the form

$$\mathcal{F}_i(x, y) = \{B_i T^t C_i : t \in \mathbb{Z}\},$$

where B_i and C_i are some matrices from $\text{SL}(2, \mathbb{Z})$. Moreover, there is a polyno-
 325 mial time algorithm that determines whether each $\mathcal{F}_i(x, y)$ is empty or not and
 in the latter case finds corresponding matrices B_i and C_i .

Proof. Let us write the numbers x and y as $x = \frac{x_1}{x_2}$ and $y = \frac{y_1}{y_2}$, where we
 assume that $\gcd(x_1, x_2) = \gcd(y_1, y_2) = 1$. Consider the equation $f_M(x) = y$,
 where $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$ is an unknown matrix from $\text{SL}(2, \mathbb{Z})$. We can rewrite it as

$$\frac{a\frac{x_1}{x_2} + b}{c\frac{x_1}{x_2} + d} = \frac{y_1}{y_2} \quad \text{or} \quad \frac{ax_1 + bx_2}{cx_1 + dx_2} = \frac{y_1}{y_2}. \quad (4)$$

330 Consider the vectors $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$, $\mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$, and $\mathbf{z} = \begin{bmatrix} z_1 \\ z_2 \end{bmatrix}$, where \mathbf{z} is the vector
 with coordinates $z_1 = ax_1 + bx_2$ and $z_2 = cx_1 + dx_2$. So we have that $\mathbf{z} = M\mathbf{x}$.
 In this notation Equation (4) is equivalent to the fact that vector $\mathbf{z} = M\mathbf{x}$
 belongs to the set $\{k\mathbf{y} : k \in \mathbb{Z}\}$.

Recall that $\gcd(x_1, x_2) = 1$ and hence, by Lemma 4, we also have that
 $\gcd(z_1, z_2) = 1$. Thus if $\mathbf{z} = k\mathbf{y}$ for some $k \in \mathbb{Z}$, then we must have that
 $k = \pm 1$. In other words, we showed that Equation (4) is equivalent to two
 matrix equations: $M\mathbf{x} = \mathbf{y}$ and $M\mathbf{x} = -\mathbf{y}$. So we have that

$$\begin{aligned} \mathcal{F}(x, y) &= \mathcal{F}_1(x, y) \cup \mathcal{F}_2(x, y), \quad \text{where} \\ \mathcal{F}_1(x, y) &= \{M \in \text{SL}(2, \mathbb{Z}) : M\mathbf{x} = \mathbf{y}\} \quad \text{and} \\ \mathcal{F}_2(x, y) &= \{M \in \text{SL}(2, \mathbb{Z}) : M\mathbf{x} = -\mathbf{y}\}. \end{aligned}$$

Note that $\mathbf{x} \neq \mathbf{0}$ because $x_2 \neq 0$. Hence by Theorem 10 and Corollary 11, each
 335 $\mathcal{F}_i(x, y)$ is either empty or has the form $\mathcal{F}_i(x, y) = \{B_i T^t C_i : t \in \mathbb{Z}\}$ for some
 B_i and C_i from $\text{SL}(2, \mathbb{Z})$ which can be computed in polynomial time. \square

5. Decidability of VRP and FLT

We now prove that the intersection emptiness problem of two regular subsets
 in $\text{SL}(2, \mathbb{Z})$ is decidable (Proposition 15). Then in the proof of Theorem 16 we
 340 will show that the solution sets of the equations $M\mathbf{x} = \mathbf{y}$ and $f_M(x) = y$ are
 regular subsets of $\text{SL}(2, \mathbb{Z})$. Using these results and the fact that the semigroups
 in $\text{SL}(2, \mathbb{Z})$ are also regular subsets, we will conclude that the vector reachability
 problem and the reachability problem by fractional linear transformations in
 $\text{SL}(2, \mathbb{Z})$ are decidable.

345 Consider an alphabet $\Sigma = \{S, R\}$ consisting of two symbols S and R and
 define the mapping $\varphi : \Sigma \rightarrow \text{SL}(2, \mathbb{Z})$ as follows:

$$\varphi(S) = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix} \text{ and } \varphi(R) = \begin{bmatrix} 0 & -1 \\ 1 & 1 \end{bmatrix}.$$

We can extend this mapping to the morphism $\varphi : \Sigma^* \rightarrow \text{SL}(2, \mathbb{Z})$ in a natural
 way. The matrices $\varphi(S)$ and $\varphi(R)$ are in fact generators of $\text{SL}(2, \mathbb{Z})$, so φ is
 surjective. We call a word $w \in \Sigma^*$ *reduced* if it does not have substrings of the
 350 form SS or RRR . In our proof we will make use of the following well-known
 fact.

Theorem 13 ([34, 35, 36]). *For every $M \in \text{SL}(2, \mathbb{Z})$, there exists a unique
 reduced word $w \in \Sigma^*$ such that either $M = \varphi(w)$ or $M = -\varphi(w)$.*

Definition 14. A *signed automaton* $\mathcal{A} = (\Sigma, Q, I, \Delta, F^+, F^-)$ is a nondeter-
 355 ministic finite automaton whose final states are divided into two (not necessarily
 disjoint) subsets F^+ and F^- .

A *signed language* accepted by a signed automaton \mathcal{A} is a pair $L(\mathcal{A}) =$
 $(L(\mathcal{A})^+, L(\mathcal{A})^-)$, where $L(\mathcal{A})^+$ and $L(\mathcal{A})^-$ consists of the words $w \in \Sigma^*$ for
 which there is a run of \mathcal{A} that ends in the set F^+ or F^- , respectively. Note
 360 that we do not assume that $L(\mathcal{A})^+$ and $L(\mathcal{A})^-$ are disjoint.

Let $L = (L^+, L^-)$ be a signed language, then we define a regular subset of
 $\text{SL}(2, \mathbb{Z})$ corresponding to this language as

$$\varphi(L) = \{\varphi(w) : w \in L^+\} \cup \{-\varphi(w) : w \in L^-\}.$$

The following proposition is an important ingredient of our main results.

Proposition 15. *There is an algorithm that for any given regular signed lan-
 365 guages L_1 and L_2 over the alphabet Σ , decides whether $\varphi(L_1) \cap \varphi(L_2)$ is empty
 or not.*

Proof. We will describe a construction that turns any signed automaton \mathcal{A} over
 Σ into a new signed automaton $\tilde{\mathcal{A}}$ such that

- $\varphi(L(\tilde{\mathcal{A}})) = \varphi(L(\mathcal{A}))$ and

- for every $M \in \varphi(L(\tilde{\mathcal{A}}))$, there is a reduced word w such that $M = \varphi(w)$ or $M = -\varphi(w)$ and $w \in L(\tilde{\mathcal{A}})^+$ or $w \in L(\tilde{\mathcal{A}})^-$, respectively.

Suppose $\mathcal{A} = (\Sigma, Q, I, \Delta, F^+, F^-)$, then define the automaton $\tilde{\mathcal{A}}$ as follows: $\tilde{\mathcal{A}} = (\Sigma, \tilde{Q}, \tilde{I}, \tilde{\Delta}, \tilde{F}^+, \tilde{F}^-)$, where

- $\tilde{Q} = Q \times \{+, -\}$,
- $\tilde{I} = I \times \{+\}$,
- $\tilde{F}^+ = \{(q, +) : q \in F^+\} \cup \{(q, -) : q \in F^-\}$,
- $\tilde{F}^- = \{(q, +) : q \in F^-\} \cup \{(q, -) : q \in F^+\}$,

To define $\tilde{\Delta}$, we first set $\tilde{\Delta} = \emptyset$. Then for each transition $(q_1, X, q_2) \in \Delta$, we add the following two transition into $\tilde{\Delta}$: $((q_1, +), X, (q_2, +))$ and $((q_1, -), X, (q_2, -))$.

Furthermore, we iteratively add new ε -transitions to $\tilde{\Delta}$ as follows: if there is a run of $\tilde{\mathcal{A}}$ from (q_1, s_1) to (q_2, s_2) labelled by SS or RRR , then we add an ε -transition from (q_1, s_1) to (q_2, \bar{s}_2) , where \bar{s}_2 is the sign opposite to s_2 . For instance, if there is a run from $(q_1, +)$ to $(q_2, +)$ labelled by RRR , then we add an ε -transition from $(q_1, +)$ to $(q_2, -)$ (see Figure 1 for an illustration). We continue this process until no new ε -transitions can be added.

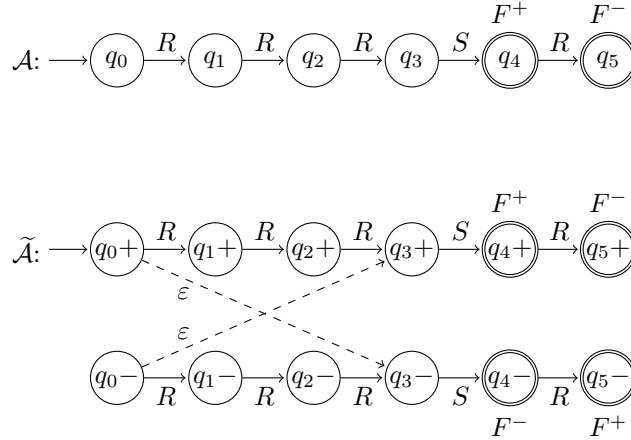


Figure 1: An example of an automaton \mathcal{A} (above) and its corresponding automaton $\tilde{\mathcal{A}}$ (below). The final states from F^+ and F^- are marked by the labels F^+ and F^- , respectively.

Note that in $\text{SL}(2, \mathbb{Z})$ we have $\varphi(S)^2 = \varphi(R)^3 = -I$, and this is reflected in the change of sign of the end state of a new ε -transition. It is not hard to see that $\tilde{\mathcal{A}}$ is indeed the desired automaton.

Let \mathcal{A}_1 and \mathcal{A}_2 be two finite signed automata such that $L(\mathcal{A}_1) = L_1$ and $L(\mathcal{A}_2) = L_2$. To check whether $\varphi(L_1) \cap \varphi(L_2)$ is empty or not, we take the

automata \mathcal{A}_1 and \mathcal{A}_2 and construct the new automata $\tilde{\mathcal{A}}_1$ and $\tilde{\mathcal{A}}_2$ as described above.

We now prove that

$$\varphi(L_1) \cap \varphi(L_2) \neq \emptyset \text{ iff } L(\tilde{\mathcal{A}}_1)^+ \cap L(\tilde{\mathcal{A}}_2)^+ \neq \emptyset \text{ or } L(\tilde{\mathcal{A}}_1)^- \cap L(\tilde{\mathcal{A}}_2)^- \neq \emptyset.$$

Indeed, suppose that $M \in \varphi(L_1) \cap \varphi(L_2)$. By the above construction we have
 395 $\varphi(L(\tilde{\mathcal{A}}_i)) = \varphi(L_i)$, for $i = 1, 2$, and there is a unique reduced word $w \in \Sigma^*$ such that $M = \varphi(w)$ or $M = -\varphi(w)$ and $w \in L(\tilde{\mathcal{A}}_i)^+$ or $w \in L(\tilde{\mathcal{A}}_i)^-$, respectively, for both $i = 1, 2$. In the first case we have $w \in L(\tilde{\mathcal{A}}_1)^+ \cap L(\tilde{\mathcal{A}}_2)^+$ and in the second case $w \in L(\tilde{\mathcal{A}}_1)^- \cap L(\tilde{\mathcal{A}}_2)^-$. The implication in the other direction is trivial.

400 To complete the proof we note that the intersection of regular languages is again regular, and the emptiness problem for regular languages is decidable. \square

We are now ready to prove our main results.

Theorem 16. *The vector reachability problem (VRP) and the reachability problem by fractional linear transformations (FLT) in $\text{SL}(2, \mathbb{Z})$ are decidable.*

405 *Proof.* Suppose M_1, \dots, M_n is a given finite collection of matrices from $\text{SL}(2, \mathbb{Z})$. Let $w_1, \dots, w_n \in \Sigma^*$ be some words, not necessarily reduced, such that $M_i = \varphi(w_i)$, for $i = 1, \dots, n$. Define the language $\mathcal{L}_{\text{semigr}}$ that corresponds to the semigroup $\langle M_1, \dots, M_n \rangle$ as $\mathcal{L}_{\text{semigr}} = (w_1 + w_2 + \dots + w_n)^*$.

Recall that in the vector reachability problem we are given two vectors \mathbf{x}
 410 and \mathbf{y} from $\mathbb{Z} \times \mathbb{Z}$, and we ask if there is a matrix $M \in \langle M_1, \dots, M_n \rangle$ such that $M\mathbf{x} = \mathbf{y}$. We want to construct a regular language $\mathcal{L}_{\mathbf{x}, \mathbf{y}}^{\text{vrp}}$ that corresponds to these matrices.

If $\mathbf{x} = \mathbf{0}$ and $\mathbf{y} \neq \mathbf{0}$, then we set $\mathcal{L}_{\mathbf{x}, \mathbf{y}}^{\text{vrp}} = \emptyset$ because in this case the equation $M\mathbf{x} = \mathbf{y}$ does not have a solution. On the other hand, if $\mathbf{x} = \mathbf{0}$ and $\mathbf{y} = \mathbf{0}$, then
 415 we set $\mathcal{L}_{\mathbf{x}, \mathbf{y}}^{\text{vrp}} = \{S, R\}^*$ because any matrix $M \in \text{SL}(2, \mathbb{Z})$ satisfies the equation $M\mathbf{0} = \mathbf{0}$.

Now assume that $\mathbf{x} \neq \mathbf{0}$. Then by Theorem 10, the matrix equation $M\mathbf{x} = \mathbf{y}$ either has no solution, or its solution has the form $\{BT^tC : t \in \mathbb{Z}\}$, where
 $T = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}$, and B and C are some matrices from $\text{SL}(2, \mathbb{Z})$. Moreover, B and
 420 C can be computed from \mathbf{x} and \mathbf{y} in PTIME. In the case when $M\mathbf{x} = \mathbf{y}$ has no solution, we set $\mathcal{L}_{\mathbf{x}, \mathbf{y}}^{\text{vrp}} = \emptyset$. If the solution set is non-empty, then we can rewrite it as

$$\{BT^tC : t \in \mathbb{Z}\} = \{BT^tC : t \geq 0\} \cup \{BT^{-t}C : t \geq 0\}.$$

Let u and v be words from Σ^* such that $B = \varphi(u)$ and $C = \varphi(v)$. It is easy to check that $T = \varphi(S^3R)$ and $T^{-1} = \varphi(R^5S)$. Hence

$$\mathcal{L}_{\mathbf{x}, \mathbf{y}}^{\text{vrp}} = u(S^3R)^*v + u(R^5S)^*v$$

425 is a regular language that describes the solutions of the equation $M\mathbf{x} = \mathbf{y}$ in $\text{SL}(2, \mathbb{Z})$.

In a similar way we can construct a regular language $\mathcal{L}_{x,y}^{\text{flt}}$ that corresponds to the reachability problem by fractional linear transformations from x to y . By Theorem 12, the set $\mathcal{F}(x, y)$ of matrices from $\text{SL}(2, \mathbb{Z})$ that satisfy the equation $f_M(x) = y$ is equal to $\mathcal{F}(x, y) = \mathcal{F}_1(x, y) \cup \mathcal{F}_2(x, y)$, where each $\mathcal{F}_i(x, y)$ is either empty or has the form $\mathcal{F}_i(x, y) = \{B_i T^t C_i : t \in \mathbb{Z}\}$, where T is as above, and B_i and C_i are some matrices from $\text{SL}(2, \mathbb{Z})$. All these matrices can be computed in PTIME from x and y .

We define $\mathcal{L}_{x,y}^{\text{flt}}$ as the union $\mathcal{L}_{x,y}^{\text{flt}} = \mathcal{L}_1 \cup \mathcal{L}_2$ of two regular languages \mathcal{L}_1 and \mathcal{L}_2 . If $\mathcal{F}_i(x, y)$ is empty, then we set $\mathcal{L}_i = \emptyset$. Otherwise, let u_i and v_i be words from Σ^* such that $B_i = \varphi(u_i)$ and $C_i = \varphi(v_i)$. Then we can define \mathcal{L}_i as

$$\mathcal{L}_i = u_i(S^3 R)^* v_i + u_i(R^5 S)^* v_i.$$

Thus we defined a regular language $\mathcal{L}_{x,y}^{\text{flt}}$ that corresponds the solution set of the equation $f_M(x) = y$ in $\text{SL}(2, \mathbb{Z})$.

We remind that in Proposition 15 we work with signed languages. Therefore, in what follows we convert every regular language L that we have constructed so far into a corresponding signed language (L, \emptyset) .

Finally, the vector reachability problem for \mathbf{x} and \mathbf{y} has a solution if and only if

$$\varphi((\mathcal{L}_{\mathbf{x},\mathbf{y}}^{\text{vfp}}, \emptyset)) \cap \varphi((\mathcal{L}_{\text{semigr}}, \emptyset)) \neq \emptyset.$$

Similarly, the reachability problem by fractional linear transformations for x and y has a solution if and only if

$$\varphi((\mathcal{L}_{x,y}^{\text{flt}}, \emptyset)) \cap \varphi((\mathcal{L}_{\text{semigr}}, \emptyset)) \neq \emptyset.$$

By Proposition 15 these questions are algorithmically decidable. \square

Remark 1. In the definition of the vector reachability problem we consider vectors \mathbf{x} and \mathbf{y} only with integer coefficients. However, the above theorem still holds if we allow \mathbf{x} and \mathbf{y} to have rational coefficients. Indeed, the equation $M\mathbf{x} = \mathbf{y}$ is equivalent to $M(\lambda\mathbf{x}) = \lambda\mathbf{y}$ for any $\lambda \neq 0$. So if we multiply the equation $M\mathbf{x} = \mathbf{y}$ by the greatest common divisor of all coefficients, we can transform it to an equivalent equation with integer coefficients.

Remark 2. A characterization of the matrices M from $\text{SL}(2, \mathbb{Z})$ that satisfy the equation $M\mathbf{x} = \mathbf{y}$, which is given in Theorem 10, can be computed in polynomial time. However the overall complexity of the algorithm is EXPTIME if the entries of the matrices are given in binary presentation. This is due to the fact that a reduced word w that corresponds to a given matrix M , i.e., such that $M = \pm\varphi(w)$, has length exponential in the binary presentation of M . So computing symbolic presentations of given matrices and constructing automata for the languages $\mathcal{L}_{\text{semigr}}$, $\mathcal{L}_{\mathbf{x},\mathbf{y}}^{\text{vfp}}$ and $\mathcal{L}_{x,y}^{\text{flt}}$ takes exponential time. The next steps of the algorithm take only polynomial time in the size of these automata. However the PTIME algorithm for computing all mappings from \mathbf{x} to \mathbf{y} could be combined with the result of Gurevich and Schupp [17] to produce a polynomial time algorithm for the vector reachability problem over the modular group.

465 In the rest of this section we will give some generalizations of the above theorem.

Consider a semigroup generated by matrices M_1, \dots, M_n from $\text{SL}(2, \mathbb{Z})$. As we showed above, this semigroup can be described by a regular language which we called $\mathcal{L}_{\text{semigr}}$. It's not hard to see that the proof of Theorem 16 remains
 470 valid if we replace $\mathcal{L}_{\text{semigr}}$ by any other regular language, that is, a language defined by a finite automaton or a labelled transition system.

Proposition 17. *Suppose that we are given a finite collection of matrices M_1, \dots, M_n from $\text{SL}(2, \mathbb{Z})$ and a regular language $L \subseteq \{1, \dots, n\}^*$. Consider the following generalized reachability problems:*

- 475 • **Generalized vector reachability problem.** *Given two vectors \mathbf{x} and \mathbf{y} with integer coefficients, decide whether there exists a word $i_1 \dots i_k$ from the language L such that $M_{i_1} \dots M_{i_k} \mathbf{x} = \mathbf{y}$.*
- **Generalized reachability problem by fractional linear transformations.** *Given two rational numbers x and y , decide whether there
 480 exists a word $i_1 \dots i_k$ from L such that $f_{M_{i_1} \dots M_{i_k}}(x) = y$.*

Then the above generalized reachability problems are decidable.

Proof. The proof of this proposition is similar to the proof of Theorem 16. Namely, it follows from the fact that a regular language L defines a regular subset in $\text{SL}(2, \mathbb{Z})$ and Proposition 15, where we proved that the emptiness
 485 problem for the intersection of two regular subsets in $\text{SL}(2, \mathbb{Z})$ is decidable. \square

As an application of Proposition 17 let us consider the follow matrix equation

$$M_1^{x_1} \dots M_k^{x_k} \mathbf{x} = N_1^{y_1} \dots N_l^{y_l} \mathbf{y}, \quad (5)$$

where x_1, \dots, x_k and y_1, \dots, y_l are non-negative integers. In [30] it was proved that if M_1, \dots, M_k and N_1, \dots, N_l are commuting $n \times n$ matrices over algebraic numbers and \mathbf{x}, \mathbf{y} are vectors with algebraic coefficients, then it is decid-
 490 able in polynomial time whether Equation (5) has a solution. On the other hand, in [37] it was shown that there is no algorithm for solving the equation $M_1^{x_1} \dots M_k^{x_k} = Z$, where M_1, \dots, M_k are integer $n \times n$ matrices and Z is the zero matrix. Using the construction of Kronecker (or tensor) product of matrices, it
 495 is possible to show that the above-mentioned result implies that Equation (5) is algorithmically undecidable in general for non-commuting integer matrices M_1, \dots, M_k and N_1, \dots, N_l .

However using Proposition 17 we can algorithmically solve Equation (5) in the case when M_1, \dots, M_k and N_1, \dots, N_l are matrices from $\text{SL}(2, \mathbb{Z})$ and the
 500 vectors \mathbf{x}, \mathbf{y} have integer coefficients. Indeed, since the matrices from $\text{SL}(2, \mathbb{Z})$ are invertible, we can rewrite (5) as

$$(N_l^{-1})^{y_l} \dots (N_1^{-1})^{y_1} M_1^{x_1} \dots M_k^{x_k} \mathbf{x} = \mathbf{y}.$$

It is not hard to see that

$$\{(N_l^{-1})^{y_l} \dots (N_1^{-1})^{y_1} M_1^{x_1} \dots M_k^{x_k} : x_1, \dots, x_k, y_1, \dots, y_l \in \mathbb{N} \cup \{0\}\}$$

is a regular subset of $\text{SL}(2, \mathbb{Z})$. Hence this problem is decidable. Using the same idea we can algorithmically solve Equation (5) also in the case when x_1, \dots, x_k and y_1, \dots, y_l are arbitrary integers and the matrices are from $\text{SL}(2, \mathbb{Z})$.

6. Geometric properties and the scalar reachability problem

In this section we will discuss a geometric interpretation of both reachability problems (Figure 2), which we will use later to prove that a special case of the scalar reachability problem is decidable (Proposition 20).

Proposition 18. *By Theorem 10, the set of matrices M from $\text{SL}(2, \mathbb{Z})$ that transform a vector $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$ to a vector $\mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$ has the form*

$$\mathcal{F} = \{BT^{kt}C : t \in \mathbb{Z}\}.$$

Consider the equation $BT^{kt}C\mathbf{x} = \mathbf{y}$ and let us make the following change of variables: $\mathbf{u} = C\mathbf{x}$ and $\mathbf{v} = B^{-1}\mathbf{y}$:

$$\mathbf{x} \xrightarrow{C} \mathbf{u} \xrightarrow{T^{kt}} \mathbf{v} \xrightarrow{B} \mathbf{y}.$$

Then $\mathbf{u} = \mathbf{v} = \begin{bmatrix} d \\ 0 \end{bmatrix}$, where $|d| = \gcd(x_1, x_2) = \gcd(y_1, y_2)$.

Proof. In the new notations, the equation $BT^{kt}C\mathbf{x} = \mathbf{y}$ can be written as $T^{kt}\mathbf{u} = \mathbf{v}$, and this equality holds for any $t \in \mathbb{Z}$. Now let $\mathbf{u} = \begin{bmatrix} u_1 \\ u_2 \end{bmatrix}$ and $\mathbf{v} = \begin{bmatrix} v_1 \\ v_2 \end{bmatrix}$. Hence we have $\begin{bmatrix} 1 & kt \\ 0 & 1 \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \end{bmatrix} = \begin{bmatrix} v_1 \\ v_2 \end{bmatrix}$, which is equivalent to $u_2 = v_2$ and $u_1 + kt u_2 = v_1$, for any $t \in \mathbb{Z}$. So, we must have $u_2 = v_2 = 0$ and hence $u_1 = v_1$.

Therefore, the vectors \mathbf{u} and \mathbf{v} have the form $\mathbf{u} = \mathbf{v} = \begin{bmatrix} d \\ 0 \end{bmatrix}$ for some $d \in \mathbb{Z}$. Moreover, since $\mathbf{u} = C\mathbf{x}$, we obtain from Lemma 4 that $|d| = \gcd(x_1, x_2) = \gcd(y_1, y_2)$. \square

Using the above proposition, we can give the following geometric interpretation of the transformation $BT^tC\mathbf{x} = \mathbf{y}$: first, we apply C to \mathbf{x} and arrive at $\mathbf{u} = \begin{bmatrix} d \\ 0 \end{bmatrix}$, then we loop at \mathbf{u} for t many times using T , and finally apply B to move from \mathbf{u} to \mathbf{y} (see Figure 2 on the left).

Similarly, we have the following geometric interpretation of the fractional linear transformation $y = f_{BT^tC}(x) = f_B \circ f_{T^t} \circ f_C(x)$: first it maps x to ∞

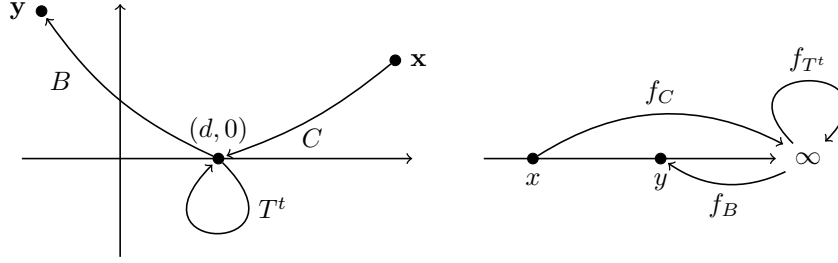


Figure 2: Geometric interpretation of the linear transformation $\mathbf{y} = BT^tC\mathbf{x}$ (left) and of the fractional linear transformation $y = f_{BT^tC}(x)$ (right).

using f_C , then loops at ∞ for t many times using f_T , and finally maps ∞ to y using f_B (see Figure 2 on the right).

We now show how to apply the geometric interpretation of the vector reachability problem to solve a special case of the scalar reachability problem.

Definition 19. The *scalar reachability problem* in $\text{SL}(2, \mathbb{Z})$ is stated as follows:

Let $[z_1, z_2]$ and $\begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$ be vectors from $\mathbb{Z} \times \mathbb{Z}$ and let λ be an integer number.

We are also given a finite collection of matrices M_1, \dots, M_n from $\text{SL}(2, \mathbb{Z})$. The question is to decide whether there exists a matrix $M \in \langle M_1, \dots, M_n \rangle$ which satisfies the equation $[z_1, z_2] M \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \lambda$.

We will consider a special case of this problem when $z_2 = 1$ and $\lambda = 1$. Our proof relies on the characterization from Theorem 10 and Corollary 11 and on Proposition 15 in which we showed that the emptiness problem for the intersection of two regular subsets in $\text{SL}(2, \mathbb{Z})$ is decidable.

Proposition 20. Suppose that the above equation has the form

$$[a, 1] M \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = 1, \quad (6)$$

where a , x_1 and x_2 are some integer numbers. Then this special case of the scalar reachability problem is decidable.

Proof. The general idea of the proof is the same as in Theorem 16, that is, we will show that the set of matrices $M \in \text{SL}(2, \mathbb{Z})$ that satisfy Equation (6) can be described by a regular language. First, let us consider a geometric interpretation of this problem. We can rewrite Equation (6) as a system of two equations: $M \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$ and $ay_1 + y_2 = 1$. So, M satisfies Equation (6) if and only if it maps a fixed vector $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$ to some vector $\mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$ that lies on the line L described by the equation $ay_1 + y_2 = 1$. In other words, we have a

vector to line reachability problem for the line L that is defined by the equation $ay_1 + y_2 = 1$.

Note that if a vector \mathbf{y} lies of the line $ay_1 + y_2 = 1$, then $\gcd(y_1, y_2) = 1$.
 555 Hence by Lemma 4, Equation (6) has a solution only if $\gcd(x_1, x_2) = 1$. So, from now on we assume that $\gcd(x_1, x_2) = 1$.

By Corollary 11, any $M \in \text{SL}(2, \mathbb{Z})$ that maps \mathbf{x} to a vector \mathbf{y} on the line L has the form $M = BT^tC$, where B and C are some matrices from $\text{SL}(2, \mathbb{Z})$ and $t \in \mathbb{Z}$. Geometrically, the transformation $\mathbf{y} = BT^tC\mathbf{x}$ goes via the point $(1, 0)$
 560 as shown in Figure 3.

Note that the matrices B and C above depend on the vector \mathbf{y} as a parameter. Here we prove a useful lemma which will imply that we can choose only one matrix C that maps \mathbf{x} to $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ independently of the vector \mathbf{y} .

Lemma 21. Let $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$ and $\mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$ be any vectors from $\mathbb{Z} \times \mathbb{Z}$ such
 565 that $\gcd(x_1, x_2) = \gcd(y_1, y_2) = d$. Let d_1 and d_2 be any integer numbers with $|d_1| = |d_2| = d$ and let A_1, B_1 and A_2, B_2 by any matrices from $\text{SL}(2, \mathbb{Z})$ such that $B_i\mathbf{x} = \begin{bmatrix} d_i \\ 0 \end{bmatrix}$ and $A_i \begin{bmatrix} d_i \\ 0 \end{bmatrix} = \mathbf{y}$, for $i = 1, 2$. Then

$$\{A_1T^tB_1 : t \in \mathbb{Z}\} = \{A_2T^tB_2 : t \in \mathbb{Z}\}.$$

In other words, the following diagrams define the same set of matrices that map \mathbf{x} to \mathbf{y} .

$$\begin{array}{ccc} \mathbf{x} & \xrightarrow{B_1} \begin{bmatrix} d_1 \\ 0 \end{bmatrix} & \xrightarrow{A_1} \mathbf{y} \\ & \uparrow \cup T^t & \\ & & \end{array} \quad \begin{array}{ccc} \mathbf{x} & \xrightarrow{B_2} \begin{bmatrix} d_2 \\ 0 \end{bmatrix} & \xrightarrow{A_2} \mathbf{y} \\ & \uparrow \cup T^t & \\ & & \end{array}$$

570

Proof. For $i = 1, 2$, define $\mathcal{F}_i = \{A_iT^tB_i : t \in \mathbb{Z}\}$. We need to show that $\mathcal{F}_1 = \mathcal{F}_2$. Suppose that $M \in \mathcal{F}_1$, that is, $M = A_1T^{t_1}B_1$ for some $t_1 \in \mathbb{Z}$. We want to show that $M \in \mathcal{F}_2$, that is, $M = A_2T^{t_2}B_2$ for some $t_2 \in \mathbb{Z}$. First, let us write M as

$$M = A_2(A_2^{-1}A_1)T^{t_1}(B_1B_2^{-1})B_2.$$

575 Note that

$$A_2^{-1}A_1 \begin{bmatrix} d_1 \\ 0 \end{bmatrix} = \begin{bmatrix} d_2 \\ 0 \end{bmatrix} \quad \text{and} \quad B_1B_2^{-1} \begin{bmatrix} d_2 \\ 0 \end{bmatrix} = \begin{bmatrix} d_1 \\ 0 \end{bmatrix}.$$

Suppose $A_2^{-1}A_1 = \begin{bmatrix} a_1 & a_2 \\ a_3 & a_4 \end{bmatrix}$ and $B_1B_2^{-1} = \begin{bmatrix} b_1 & b_2 \\ b_3 & b_4 \end{bmatrix}$. Then we have $a_1d_1 = d_2$, $a_3d_1 = 0$ and $b_1d_2 = d_1$, $b_3d_2 = 0$. Hence $a_3 = b_3 = 0$.

Now we need to consider two cases: (1) $d_1 = d_2$ and (2) $d_1 = -d_2$. In the first case $a_1 = b_1 = 1$ and in the second case $a_1 = b_1 = -1$. Note that $A_2^{-1}A_1$

580 and $B_1 B_2^{-1}$ are matrices from $\text{SL}(2, \mathbb{Z})$. Therefore, in the first case we must have that

$$A_2^{-1} A_1 = \begin{bmatrix} 1 & a_2 \\ 0 & 1 \end{bmatrix} \text{ and } B_1 B_2^{-1} = \begin{bmatrix} 1 & b_2 \\ 0 & 1 \end{bmatrix}$$

and in the second case we must have that

$$A_2^{-1} A_1 = \begin{bmatrix} -1 & a_2 \\ 0 & -1 \end{bmatrix} \text{ and } B_1 B_2^{-1} = \begin{bmatrix} -1 & b_2 \\ 0 & -1 \end{bmatrix}.$$

So, in the first case we obtain

$$M = A_2 T^{a_2} T^{t_1} T^{b_2} B_2 = A_2 T^{a_2+t_1+b_2} B_2,$$

and in the second case we have

$$M = A_2 (-T^{-a_2}) T^{t_1} (-T^{-b_2}) B_2 = A_2 T^{-a_2+t_1-b_2} B_2.$$

585 Hence in both cases M is in \mathcal{F}_2 . Similarly, we can show that if $M \in \mathcal{F}_2$, then $M \in \mathcal{F}_1$. \square

By Lemma 21, we can choose any matrix C from $\text{SL}(2, \mathbb{Z})$ that maps a vector \mathbf{x} to the vector $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$, and for each \mathbf{y} on the line L we can choose any matrix $B_{\mathbf{y}}$ that maps $\begin{bmatrix} 1 \\ 0 \end{bmatrix}$ to the vector \mathbf{y} . Then the solution of Equation (6)
590 will be described by the following set

$$\mathcal{F} = \{B_{\mathbf{y}} T^t C : \mathbf{y} \in L \text{ and } t \in \mathbb{Z}\}.$$

Figure 3 gives geometric interpretation of this solution.

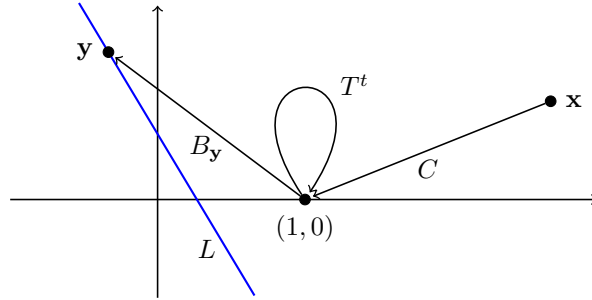


Figure 3: Geometric interpretation of the scalar reachability problem.

We need to choose $B_{\mathbf{y}}$ in such a way that the set \mathcal{F} becomes regular. If $\mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$ is a vector on the line L , then we have $ay_1 + y_2 = 1$. As one can check, if we let $B_{\mathbf{y}} = \begin{bmatrix} y_1 & -1 \\ -ay_1 + 1 & a \end{bmatrix}$ then $B_{\mathbf{y}} \in \text{SL}(2, \mathbb{Z})$ and $B_{\mathbf{y}} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \mathbf{y}$.

595 Since every entry of $B_{\mathbf{y}}$ is a linear function of y_1 , we obtain by Proposition 9
that $B_{\mathbf{y}} = AT^{ky_1}D$, where A and D are some matrices from $\text{SL}(2, \mathbb{Z})$ and k is
some integer number (in fact, one can show that $k = 1$). Finally, we can write
all solutions of Equation (6) as

$$\mathcal{F} = \{AT^{ky_1}DT^tC : y_1 \in \mathbb{Z} \text{ and } t \in \mathbb{Z}\}.$$

This is clearly a regular set and, therefore, the scalar reachability problem is
600 decidable. \square

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